

INFINITE SPECTRA IN THE FIRST ORDER THEORY OF GRAPHS

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The psectrum $\text{Spec}(A)$ of a sentence A is, roughly, the set of those α for which A has a threshold function at or near $p=n^{-\alpha}$. Examples are given of A with infinite spectra and with spectra of order type ω^i for arbitrary i .

In their fundamental work *On the Evolution of Random Graphs* Paul Erdős and Alfréd Rényi [1] showed that many natural graphtheoretic properties A possessed a threshold function $p(n)$, that is, a function having the property the if $r(n) \ll p(n)$ the random graph $G(n, r(n))$ a.s. did not satisfy A while if $p(n) \ll r(n)$ it a.s. did satisfy A . With this as motivation with S. Shelah [2] we defined the spectrum $\text{Spec}(A)$ to be, roughly, those α for which there is a threshold function near $n^{-\alpha}$. Precisely, $\alpha \in \text{Spec}(A)$ if there is a positive ε and δ either zero or one so that for any $p(n)$ satisfying $n^{-\alpha-\varepsilon} < p(n) < n^{-\alpha-\delta}$ the probability that A holds in $G(n, p(n))$ tends to δ . For our purposes we note that if for all sufficiently small ε A holds a.s. in $G(n, n^{-\alpha+\varepsilon})$ and $\neg A$ a.s. in $G(n, n^{-\alpha-\varepsilon})$, or the reverse, then $\alpha \in \text{Spec}(A)$.

We restrict our attention to sentences A of the First Order Theory of Graphs. This language contains all Boolean connectives ($\wedge, \vee, \neg, \dots$), an infinite sequence of variables x, y, z, \dots , existential and universal quantification ($\exists x$), ($\forall x$) and the predicates equality ($x=y$) and adjacency ($x \sim y$). As examples, we may express: There are no isolated points:

$$(\forall x)(\exists y)(x \sim y).$$

There is a triangle:

$$(\exists x)(\exists y)(\exists z)[x \sim y \wedge x \sim z \wedge y \sim z].$$

Every point lies in an edge not in a triangle:

$$(\forall x)(\exists y)[x \sim y] \wedge (\exists z)[z \sim x \wedge z \sim y].$$

However, many basic graphtheoretic properties such as connectivity, planarity and Hamiltonicity cannot be expressed in this language.

When A has the form "There is a copy of H " for a fixed graph H (where here copy is not necessarily an induced copy) then $\text{Spec}(A)$ was found by Erdős and Rényi.

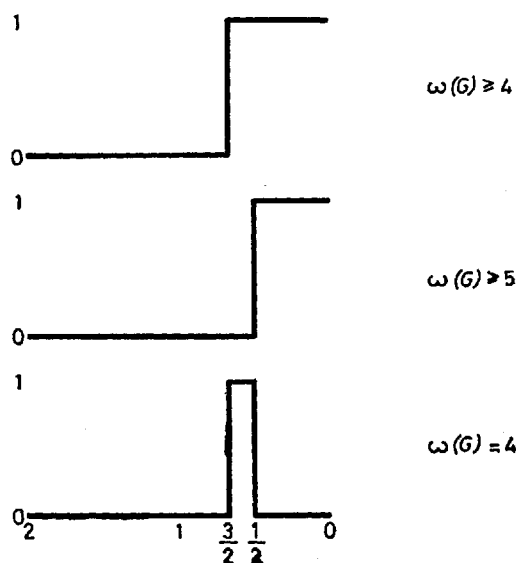


Fig. 1

In a nutshell, $\text{Spec}(A) = \{\alpha\}$ where $\alpha = v/e$ and v, e are the number of vertices and edges respectively of H unless H has a subgraph with a lower such ratio in which case $\alpha = v'/e'$ where H' with v' vertices, e' edges is the subgraph with the lowest such ratio. Letting, for example, A be "There is a K_4 " and B be "There is a K_5 ", $\text{Spec}(A) = \{2/3\}$ and $\text{Spec}(B) = \{1/2\}$. Indeed, for all rational α , $0 < \alpha \leq 1$ there is an A of this form with $\text{Spec}(A) = \{\alpha\}$. When A, B have disjoint spectra $\text{Spec}(A \vee B) = \text{Spec}(A) \cup \text{Spec}(B)$. This is illustrated in Figure 1 where the probability of $G(n, n^{-\alpha})$ having various properties is graphed versus α and the origins of the term "spectrum" become apparent. By taking a finite sequence of exclusive ors $C = A_1 \vee \dots \vee A_n$ we may find explicit first order C having as spectra any desired finite set $\{\alpha_1, \dots, \alpha_n\}$ of rational numbers in $(0, 1]$.

In [1] we showed that $\text{Spec}(A)$ consists only of rational numbers and is scattered of finite rank. Using those methods we further may show that $\text{Spec}(A)$ is well-ordered under $>$. That is, $\text{Spec}(A)$ must be a well ordered set of rational numbers of order type less than ω^ω . (Note that as $p = n^{-\alpha} >$ " is the "natural" ordering from empty to full.)

From the examples above it is not clear that $\text{Spec}(A)$ could be an infinite set. In [2] an example was given of an A with infinite spectrum. The interpretation of that A gave an arithmetization of a fragment of the positive integers. In section 1. we give a simpler example which uses only the notion of parity. In section 2. we show that for all i there exists an A whose spectrum has order type at least ω^i . It is our hope that these methods will be extended to give a full characterization of the possible spectra of first order sentences.

1. An Infinite Spectrum

Here we give a sentence A with $\text{Spec}(A) = \overline{\left\{\frac{1}{3} + \frac{1}{k}\right\}}$. Set

$$N(x, y, z) = \{u: u \sim x \wedge u \sim y \wedge u \sim z\}$$

$$H(x, y, z): N(x, y, z) \neq \emptyset$$

$$H^w(x, y): H(x, y, w)$$

Let $p = n^{-1/3-\varepsilon}$ with $\frac{1}{k+1} < \varepsilon < \frac{1}{k}$. The threshold for existence of the complete bipartite $K_{3,t}$ is when $n^{3+t}p^{3t} = 1$ so for the given p a.s. the maximal $|N(x, y, z)|$ in $G(n, p)$ is k . We shall create an A with the interpretation that the maximal $|N(x, y, z)|$ is even. Note that for any fixed $x, y, w \in V(G)$, $\Pr[H^w(x, y)] \sim np^3 \sim n^{-3\varepsilon}$.

Lemma. Let $\{x_i, y_i\}$, $1 \leq i \leq s$ be distinct pairs of elements of $V(G)$, though possibly with overlapping vertices. Assume $s \leq k/3$. Then the number of w with $H^w(x_i, y_i)$, $1 \leq i \leq s$ is a.s. $n^{1-3\varepsilon s+O(1)}$. Moreover, a.s. for all choices of $\{x_i, y_i\}$ the number of such w is $n^{1-3\varepsilon s+O(1)}$.

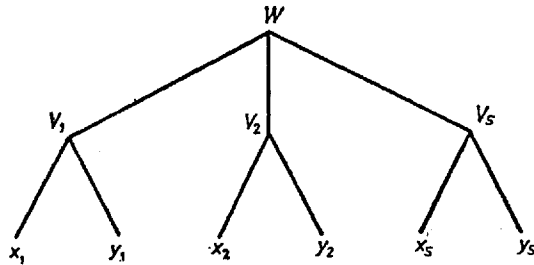


Fig. 2

Proof. In the notation of [2] the rooted graph with roots $x_1, y_1, \dots, x_s, y_s$ and non-roots w, v_1, \dots, v_s with v_i adjacent to w, x_i, y_i is a hinged extension for $\alpha = \frac{1}{3} + \varepsilon$ and so this result follows from Theorem 3 and Lemma 4 of [2]. ■

Theorem (Universal Property). Let $S \subset V(G)$, $|S| = 50k$ and let H be a graph on S with $s \leq k/3$ edges. Then there are $n^{1-3\varepsilon s+O(1)}$ w so that H^w on S is H .

Proof. The lemma gives that there are $n^{1-3\varepsilon s+O(1)}$ w so that H^w on S contains H . For any of the $O(1)$ graphs H' on S consisting of H and one more edge there are only $n^{1-3\varepsilon(s+1)+O(1)} = o(n^{1-3\varepsilon s+O(1)})$ w with H^w containing H' , deleting these give the desired w . ■

Remark. $50k$ may be replaced by any c_k .

Corollary (Extended Universal Property). Let $S \subset V(G)$, $|S| = 50k$. Let H be a graph on S with $s \leq 10k$ edges. Then there exist w_1, \dots, w_{30} so that, on S , H is the union of the H^{w_i} .

Proof. Partition H into $H_1 \dots H_{30}$, each with at most $k/3$ edges and for each H_i let w_i give H_i by the Theorem.

Notation. For any set $W = \{w_1, \dots, w_c\}$ we write $H^W(x, y)$ as shorthand for $H^{w_1}(x, y) \vee \dots \vee H^{w_c}(x, y)$. We write $(\exists_w |W| = c) \dots$ as shorthand for $(\exists_{w_1 w_2 \dots w_c}) \dots$. This notation can only be employed when c is a constant.

Now let S, T be sets whose membership is given by first order statements. We write

$$\text{BIGGER}(S, T): (\exists_w |w| = 30) (\forall_{x \in T-S} \exists !_{y \in S-T} H^w(x, y) \wedge \\ \wedge (\forall_{x, x' \in T-S} \forall_{y \in S-T} H^w(x, y) \wedge H^w(x', y) \Rightarrow x = x')).$$

If $\text{BIGGER}(S, T)$ then H^w defines an injection from $T-S$ to $S-T$ and so $|T-S| \leq |S-T|$ and thus $|S| \geq |T|$. Crucial is the partial converse. Suppose $|T| \leq |S| \leq 5k$. Let $f: T-S \rightarrow S-T$ be an injection and let $H = \{(x, f(x)): x \in T-S\}$. Then H has at most $5k$ edges and by the Extended Universal Property there is a W so that H^W is H on $S \cup T$, hence $\text{BIGGER}(S, T)$. To summarize: If $|S|, |T| \leq 5k$ then $\text{BIGGER}(S, T)$ if and only if S is bigger (or equal) to T .

Similarly we write

$$\text{EVEN}(S): (\exists_w |w| = 30) \forall_{x \in S} \forall \exists !_{y \in S} y \neq x \wedge H^w(x, y).$$

If $\text{EVEN}(S)$, H^w defines a matching on S so $|S|$ is even. If $|S| \leq 5k$ and $|S|$ is even let H be any matching on S . By the Extended Universal Property there is a W with $H^W = H$ so $\text{EVEN}(S)$. To summarize: When $|S| \leq 5k$ $\text{EVEN}(S)$ if and only if S has even size. Now write

$$\text{MAX}(x, y, z): (x')(y')(z') \text{BIGGER}[N(x, y, z), N(x', y', z')]$$

In $G(n, p)$ a.s. $\max |N(x, y, z)| = k$. Thus BIGGER is bigger (or equal) and $\text{MAX}(x, y, z)$ holds if and only if $|N(x, y, z)| = k$. Now we give our sentence.

$$A: (\exists_{x, y, z}) [\text{MAX}(x, y, z) \wedge \text{EVEN}(N(x, y, z))].$$

As EVEN is even in this range A holds a.s. when k is even and $\neg A$ holds a.s. when k is odd so $\frac{1}{3} + \frac{1}{k} \in \text{Spec}(A)$ and, as spectra are always closed, $\frac{1}{3} \in \text{Spec}(A)$. We omit the argument, it not being critical, that these are precisely the points of $\text{Spec}(A)$.

2. Spectra of Order ω^i

In this section we find by induction a sentence A_i with spectrum of order type at least ω^i . For convenience of notation we drop the index i from the predicates and sets A , EXT , UNIV , N , BIGGER , MAX , EVEN to be defined in this section. We define a critical 4-ary predicate $\text{EXT}[x, y, z, u]$. (For $i=1$ in section 1. this would be $u \sim x \wedge u \sim y \wedge u \sim z$.) For convenience we use auxilliary predicate

$$\text{UNIV}[x, y, z]: (\exists u) \text{EXT}[x, y, z, u]$$

and set

$$N[x, y, z] = \{u: \text{EXT}[x, y, z, u]\}.$$

Let $H((V(H), E(H)))$ be a graph on a subset of $[n]$, the vertex set of the random graph $G(n, p)$. (H is not necessarily a subgraph of G .) We say w represents H if for all $x, y \in V(H)$

$$\{x, y\} \in E(H) \Leftrightarrow \text{UNIV}[x, y, w].$$

Let $3 \ll a_1 \ll a_2 \ll \dots \ll a_i$ and let ε be sufficiently small, where " \ll " and "sufficiently small" may be explicitly defined. Let

$$p = n^\alpha, \quad \alpha = \frac{1}{3} + \frac{2}{9a_1} + \frac{2}{9a_1a_2} + \dots + \frac{2}{9a_1 \dots a_i} (1 + \varepsilon).$$

Our induction hypothesis (on i) is that in $G(n, p)$ a.s.

(1) All H with at most $50a_i$ vertices and at most $\frac{1}{2}a_i$ edges are represented by

some w .

(2) $\max |N[x, y, z]| = 3a_i - 1$.

(3) For all t , $0 \leq t \leq 3a_i - 1$ there are arbitrarily many disjoint x, y, z with $|N[x, y, z]| = t$.

Suppose these properties hold for i . We first give a convenient technical extension of (1). We say $W = \{w_1, \dots, w_{20}\}$ represents $H = (V(H), E(H))$ if for all $x, y \in V(H)$, $\{x, y\} \in E(H) \Leftrightarrow \text{UNIV}[x, y, w_1] \vee \dots \vee \text{UNIV}[x, y, w_{20}]$. Let H be any graph on at most $50a_i$ vertices and at most $10a_i$ edges. Decompose H into twenty graphs, each with at most $\frac{1}{2}a_i$ edges, and for each find, by (1), a w representing it.

Then W will represent H . That is,

(1') All H with at most $50a_i$ vertices and at most $10a_i$ edges are represented by some 20-set W .

We write $(\exists W)$ as shorthand for $(\exists w_1, \dots, w_{20})$ and UNIV^W for the graph it represents. Define

BIGGER $[S, T]$: $(\exists W) \text{UNIV}^W$ gives an injection from $T - S$ to $S - T$.

Then when $|S|, |T| \leq 5a_i$, **BIGGER** $[S, T] \Leftrightarrow |S| \geq |T|$. Define

MAX $[x, y, z]$: $(x')(y')(z')$ **BIGGER** $[N[x, y, z], N[x', y', z']]$.

From (2), $\max |N[x, y, z]| = 3a_i - 1$ so **BIGGER** is bigger (or equal) on these sets and thus **MAX** $[x, y, z] \Leftrightarrow |N[x, y, z]| = 3a_i - 1$. Define

EVEN $[S]$: $(\exists W) \text{UNIV}^W$ is a matching on S .

Then for sets of size at most $20a_i$, **EVEN** $[S] \Leftrightarrow |S|$ is even. Now define the sentence

A : $(\exists x, y, z) \text{MAX}[x, y, z] \wedge \text{EVEN}[N[x, y, z]]$.

When a_i is odd $3a_i - 1$ is even and a.s. A . When a_i is even, $3a_i - 1$ is odd and a.s. $\neg A$. As this holds for all ε sufficiently small there must be an $\alpha \in \text{Spec}(A)$ with

$$\frac{1}{3} + \frac{2}{9a_1} + \dots + \frac{2}{9a_1 \dots a_i} \cong \alpha \cong \frac{1}{3} + \frac{2}{9a_1} + \dots + \frac{2}{9a_1 \dots a_{i-1}(a_i + 1)}$$

and so $\text{Spec}(A)$ must have order type at least ω^i .

We now use UNIV^W to express various sizes of sets. A set S has a_i elements if there exist x, y, z so that $\text{MAX}[x, y, z]$ and $\text{MAX}[x, y, z] \cap S = \emptyset$ and there exists W so that UNIV^W on $S \times N[x, y, z]$ has degree 3 for all vertices of S except one with degree 2 and degree 1 for all vertices of $N[x, y, z]$. As we are using strong induction we may also express that S has a_j elements for any $j \leq i$. We say that S has $\frac{1}{2} a_i$ elements if there exist x, y, z so that $N[x, y, z]$ has a_i elements (described above), is disjoint from S , and there is a W so that UNIV^W on $S \times N[x, y, z]$ has degree 2 for all elements of S and degree 1 for all elements of $N[x, y, z]$. A set S has $a_i - 1$ elements if there exists $x \notin S$ so that $S \cup \{x\}$ has a_i elements. The property that a set has $a_i - c$ elements can be similarly expressed for any constant c .

We now define $\text{EXT}_{i+1}[x, y, z, u]$. For notational convenience we let EXT^* , UNIV^* , N^* denote these sets and predicates with index $i+1$. The predicate $\text{UNIV}^*[x, y, z]$ will be an extension statement with

$$v = \frac{3}{2} a_i \dots a_1 + a_i \dots a_2 + a_i \dots a_3 + \dots + a_i + 1$$

vertices (excluding x, y, z) and

$$e = \frac{9}{2} a_i \dots a_1$$

edges. We illustrate the case $i+1=4$ in Figure 3.

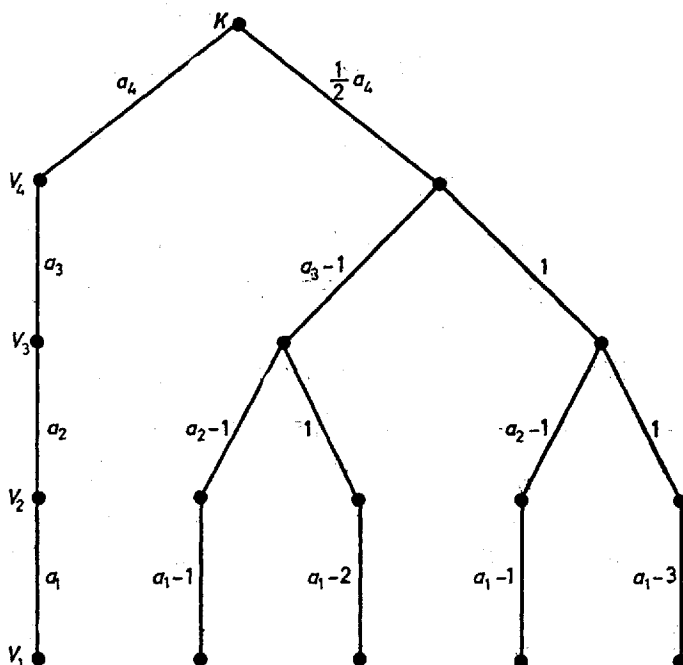


Fig. 3

The values on the edges indicate the degree of the upper level to the lower level. The bottom nodes all are joined to two of x, y, z . We read off the tree to give UNIV^* . Explicitly (letting L , left, and R , right, serve as markers for where we are)

$\text{UNIV}_4[x, y, z]$: There exists u so that

(L) There exist $a_4 v_4$'s which are adjacent to $a_3 v_3$'s which are adjacent to $a_2 v_2$'s which are adjacent to $a_1 v_1$'s which are adjacent to two of x, y, z and

(R) There exist $\frac{1}{2} a_4 v_4$'s which

(RL) are adjacent to $a_3 - 1 v_3$'s which

(RLL) are adjacent to $a_2 - 1 v_2$'s which are adjacent to $a_1 - 1 v_1$'s which are adjacent to two of x, y, z

and

(RLR) are adjacent to one v_3 which is adjacent to $a_1 - 2 v_1$'s which are adjacent to two of x, y, z

and

(RR) is adjacent to one v_3 which

(RRL) is adjacent to $a_2 - 1 v_2$'s which are adjacent to $a_1 - 1 v_1$'s adjacent to two of x, y, z

and

(RRR) is adjacent to one v_3 which is adjacent to $a_1 - 3 v_1$'s adjacent to two of x, y, z .

We omit the general case, which follows the same pattern. The values v, e are given by a straightforward calculation.

Let $a_{i+1} \gg a_i$, let ε be sufficiently small, and set

$$\alpha = \frac{1}{3} + \frac{2}{9a_1} + \dots + \frac{2}{9a_1 \dots a_i} + \frac{2(1+\varepsilon)}{9a_1 \dots a_{i+1}}.$$

Note that

$$n^v p^e = n^{o-\alpha e} = n * * \left[\frac{-1}{a_{i+1}} (1+\varepsilon) \right].$$

We first show (1) for UNIV^* and $i+1$. Let $q \leq a_{i+1}/2$ (actually, $q < a_{i+1}$ would do) and let $x_1, y_1, \dots, x_q, y_q$ be distinct pairs. The number of w with $\text{UNIV}^*[x_s, y_s, w]$, $1 \leq s \leq q$ is the number of extensions of x_1, \dots, y_q by an extension with $1+qv$ vertices and qe edges. The extension is hinged so the number of w is within a constant factor of the number of extensions which is asymptotically

$$n^{1+qv} p^{qe} = n * * \left[1 - \frac{7}{a_{i+1}} (1+\varepsilon) \right].$$

For any H on $50a_{i+1}$ vertices and q edges there are this many w with UNIV^w containing H . For each H' consisting of H and one more edge the number of w with UNIV^w containing H' has a smaller exponent and so is negligible, even when multiplied by the $O(1)$ possible H' . Thus for asymptotically this many w UNIV^w represents H , giving (1).

The number of x, y, z with $|Nc(x, y, z)| \cong s$ is the number of copies of a graph with $3+sv$ vertices and se edges. This graph is balanced. The number of copies is asymptotically $n^{3+sv} p^{se} = n * * \left[3 - \frac{s}{a_{i+1}} (1 + \varepsilon) \right]$. As the exponent decreases in s for each s there will be asymptotically this many x, y, z with $|N^*(x, y, z)| = s$, giving (3). Choosing ε sufficiently small the exponent will be positive for $s = 3a_{i+1} - 1$ so there will be x, y, z with $|N^*(x, y, z)| = 3a_{i+1} - 1$ but for arbitrarily small ε the exponent is negative for $s = 3a_{i+1}$ and so such graphs a.s. do not exist. This shows (2), completing the induction.

References

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